

BITS, BYTES, AND INTEGERS

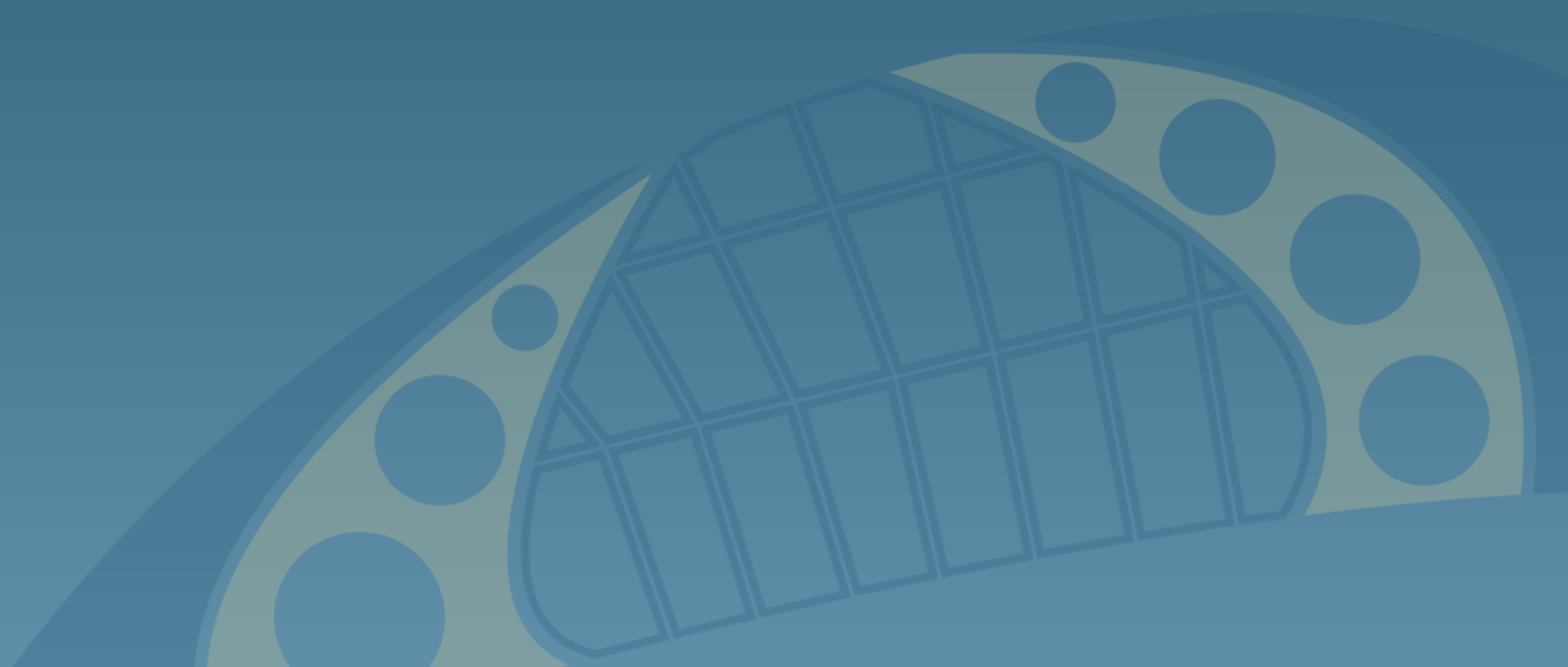
CS 045

Computer Organization and
Architecture

Prof. Donald J. Patterson

Adapted from Bryant and O'Hallaron,
Computer Systems:
A Programmer's Perspective, Third Edition

WARM-UP



WORK IT OUT



WORK IT OUT

$$T2U_4(x)$$

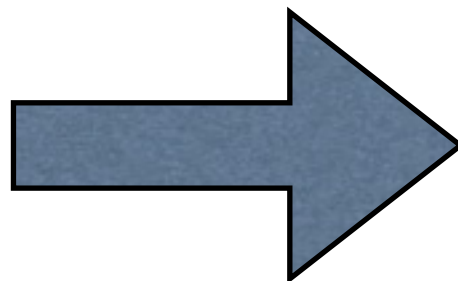
-8
-3
-2
-1
0
5



WORK IT OUT

$$T2U_4(x)$$

-8
-3
-2
-1
0
5



?
?
?
?
?
?



WORK IT OUT

$$T2U_4(x)$$

-8
-3
-2
-1
0
5

?
?
?
?
?
?



WORK IT OUT

$$T2U_4(x)$$

-8

1000₂

-3

-2

-1

0

5



WORK IT OUT

$$T2U_4(x)$$

-8

1000₂

8

-3

?

-2

?

-1

?

0

?

5

?



WORK IT OUT

$$T2U_4(x)$$

-8

1000₂

8

-3

1101₂

-2

?

-1

?

0

?

5

?



WORK IT OUT

$$T2U_4(x)$$

-8

1000₂

8

-3

1101₂

13

-2

?

-1

?

0

?

5

?



WORK IT OUT

$$T2U_4(x)$$

-8

1000₂

8

-3

1101₂

13

-2

1110₂

-1

?

0

?

5

?



WORK IT OUT

$$T2U_4(x)$$

-8

1000₂

8

-3

1101₂

13

-2

1110₂

14

-1

?

0

?

5

?



WORK IT OUT

$$T2U_4(x)$$

-8	1000 ₂	8
-3	1101 ₂	13
-2	1110 ₂	14
-1	1111 ₂	
0		?
5		?



WORK IT OUT

$$T2U_4(x)$$

-8	1000 ₂	8
-3	1101 ₂	13
-2	1110 ₂	14
-1	1111 ₂	15
0		?
5		?



WORK IT OUT

$$T2U_4(x)$$

-8	1000_2	8
-3	1101_2	13
-2	1110_2	14
-1	1111_2	15
0	0000_2	
5		?



WORK IT OUT

$$T2U_4(x)$$

-8	1000_2	8
-3	1101_2	13
-2	1110_2	14
-1	1111_2	15
0	0000_2	0
5		?



WORK IT OUT

$$T2U_4(x)$$

-8	1000_2	8
-3	1101_2	13
-2	1110_2	14
-1	1111_2	15
0	0000_2	0
5	0101_2	



WORK IT OUT

$$T2U_4(x)$$

-8	1000_2	8
-3	1101_2	13
-2	1110_2	14
-1	1111_2	15
0	0000_2	0
5	0101_2	5



WORK IT OUT

```
#include <stdio.h>

int main(){

    char a = -8;
    char b = -3;
    char c = -2;
    char d = -1;
    char e = 0;
    char f = 5;
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", a, a&0x0F, a&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", b, b&0x0F, b&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", c, c&0x0F, c&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", d, d&0x0F, d&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", e, e&0x0F, e&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", f, f&0x0F, f&0x0F);

    return 0;
}
```



WORK IT OUT

```
#include <stdio.h>

int main(){

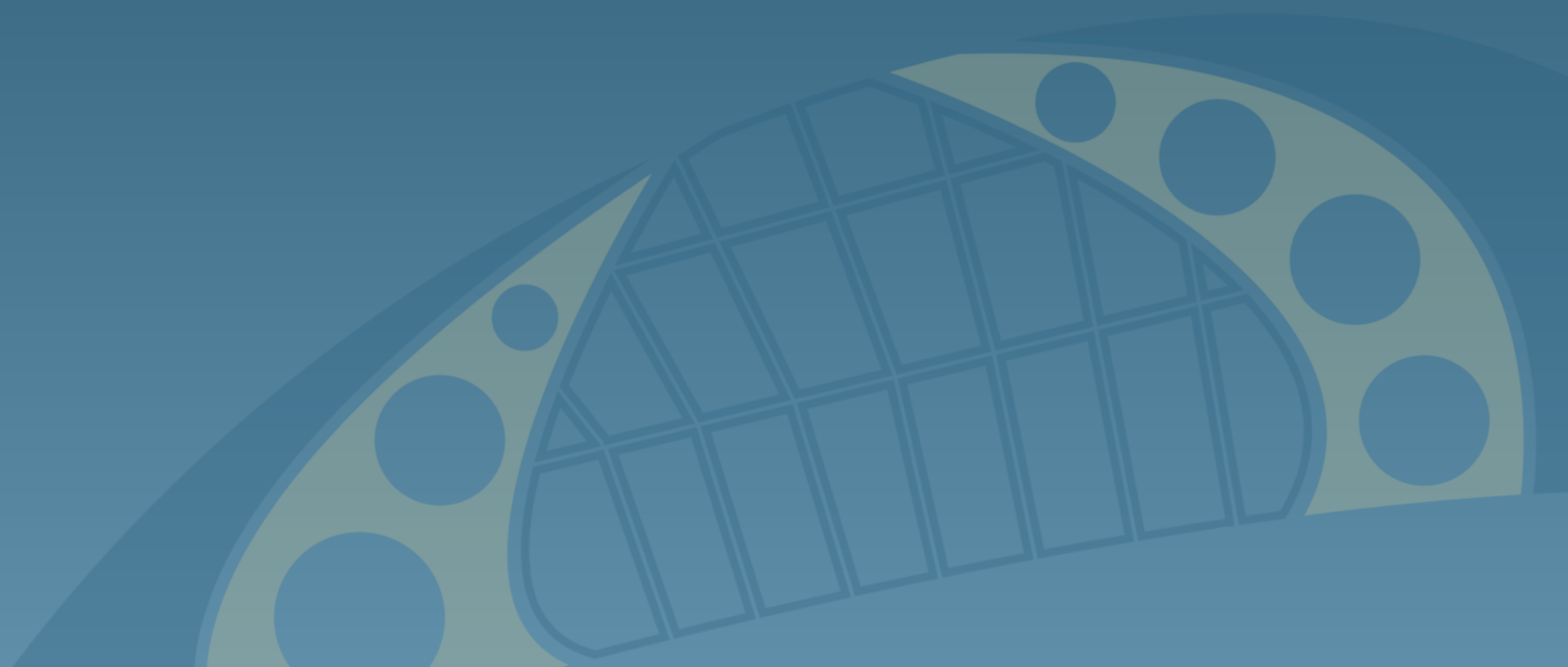
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    char d = -1;
    char e = 0;
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    printf("%2d signed -> 0x%02x -> %2u unsigned\n", a, a&0x0F, a&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", b, b&0x0F, b&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", c, c&0x0F, c&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", d, d&0x0F, d&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", e, e&0x0F, e&0x0F);
    printf("%2d signed -> 0x%02x -> %2u unsigned\n", f, f&0x0F, f&0x0F);

    return 0;
}
```

-8	signed	->	0x08	->	8	unsigned
-3	signed	->	0x0d	->	13	unsigned
-2	signed	->	0x0e	->	14	unsigned
-1	signed	->	0x0f	->	15	unsigned
0	signed	->	0x00	->	0	unsigned
5	signed	->	0x05	->	5	unsigned



EXPANDING AND TRUNCATING



CONVERSION AND CASTING

SIGN EXTENSION

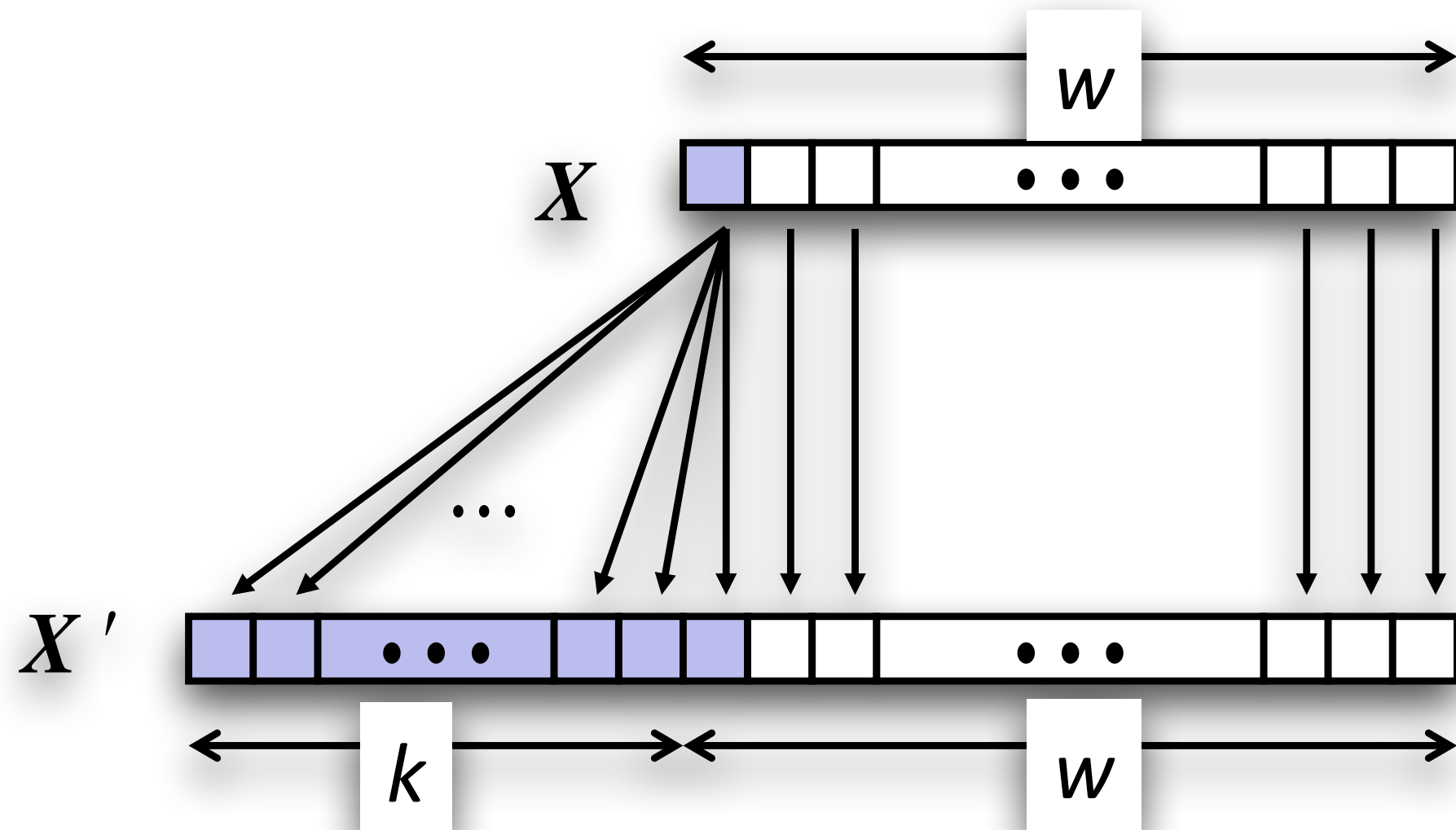
- Task:
 - Given w -bit signed integer x
 - Convert it to $w+k$ -bit integer with same value
- Rule:
 - Make k copies of sign bit:
 - $X' = \underbrace{X_{w-1}, \dots, X_{w-1}}_{k \text{ copies of MSB}}, X_{w-1}, X_{w-2}, \dots, X_0$

k copies of MSB



CONVERSION AND CASTING

SIGN EXTENSION



CONVERSION AND CASTING

SIGN EXTENSION EXAMPLE

```
short int x = 15213;  
int      ix = (int) x;  
short int y = -15213;  
int      iy = (int) y;
```

	Decimal	Hex	Binary
x	15213	3B 6D	00111011 01101101
ix	15213	00 00 3B 6D	00000000 00000000 00111011 01101101
y	-15213	C4 93	11000100 10010011
iy	-15213	FF FF C4 93	11111111 11111111 11000100 10010011

- Converting from smaller to larger integer data type
- C automatically performs sign extension



CONVERSION AND CASTING

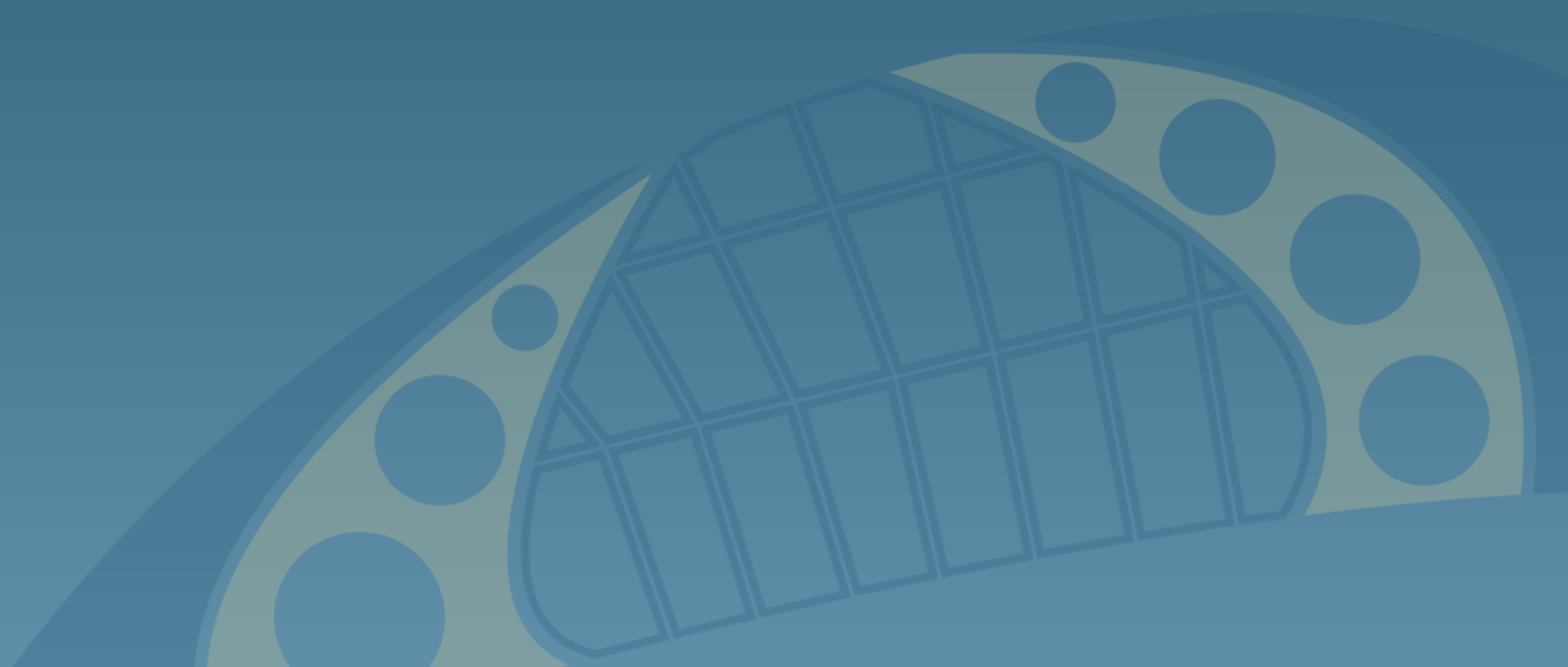
SUMMARY:

EXPANDING, TRUNCATING: BASIC RULES

- Expanding (e.g., short int to int)
 - Unsigned: zeros added
 - Signed: sign extension
 - Both yield expected result
- Truncating (e.g., unsigned to unsigned short)
 - Unsigned/signed: bits are truncated
 - Result reinterpreted
 - Unsigned: mod operation
 - Signed: similar to mod
 - For small numbers yields expected behavior



INTEGER ARITHMETIC



INTEGER ARITHMETIC

UNSIGNED ADDITION

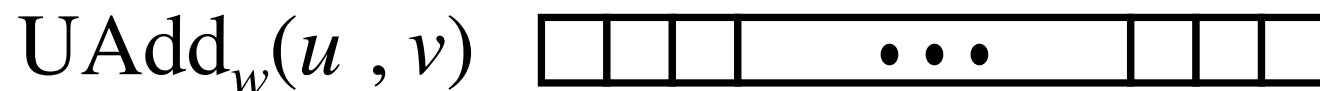
Operands: w bits



True Sum: $w+1$ bits



Discard Carry: w bits



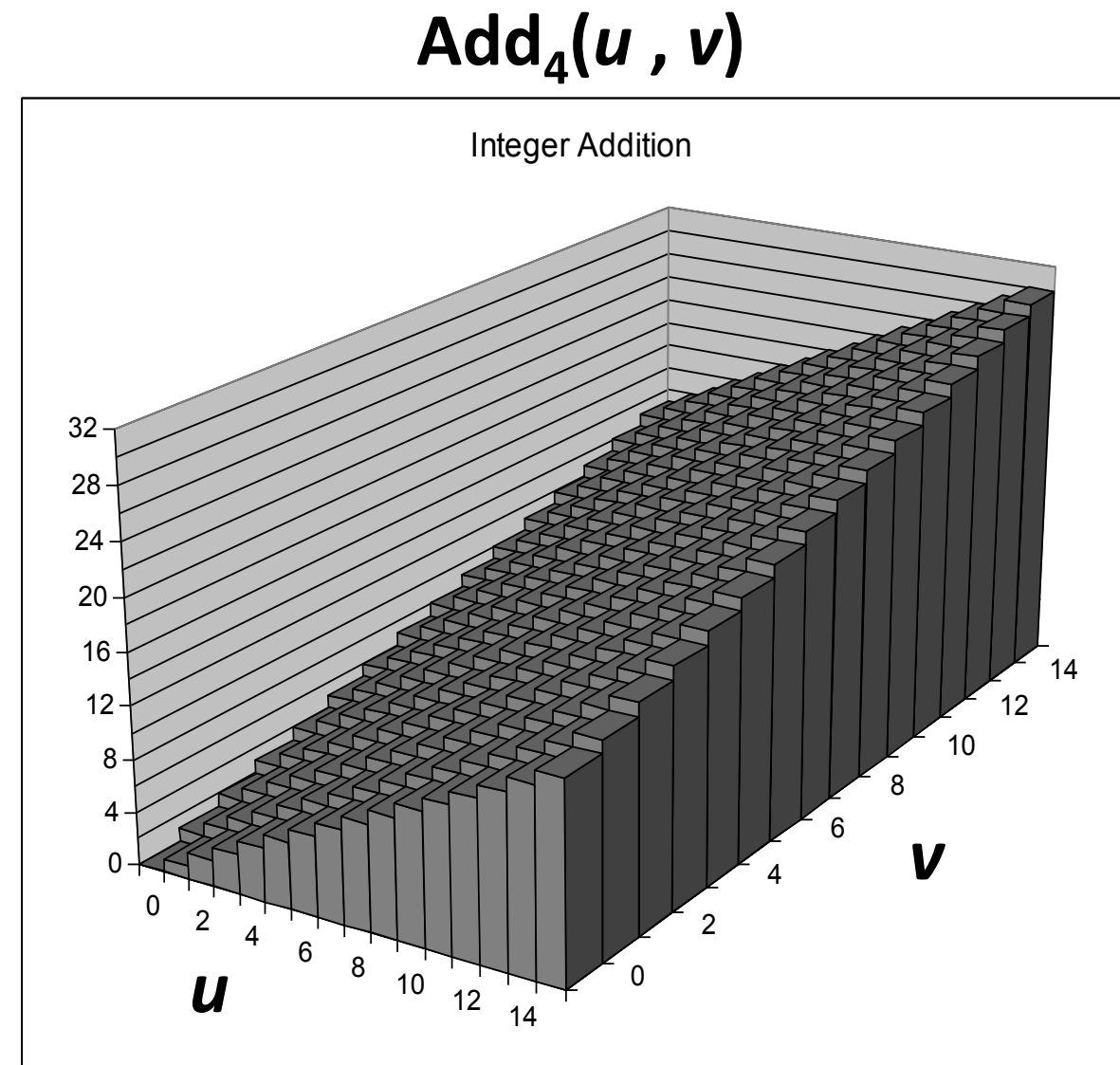
- Standard Addition Function
 - Ignores carry output
- Implements Modular Arithmetic
 - $s = \text{UAdd}_w(u, v) = (u + v) \bmod 2^w$



INTEGER ARITHMETIC

VISUALIZING (MATHEMATICAL) INTEGER ADDITION

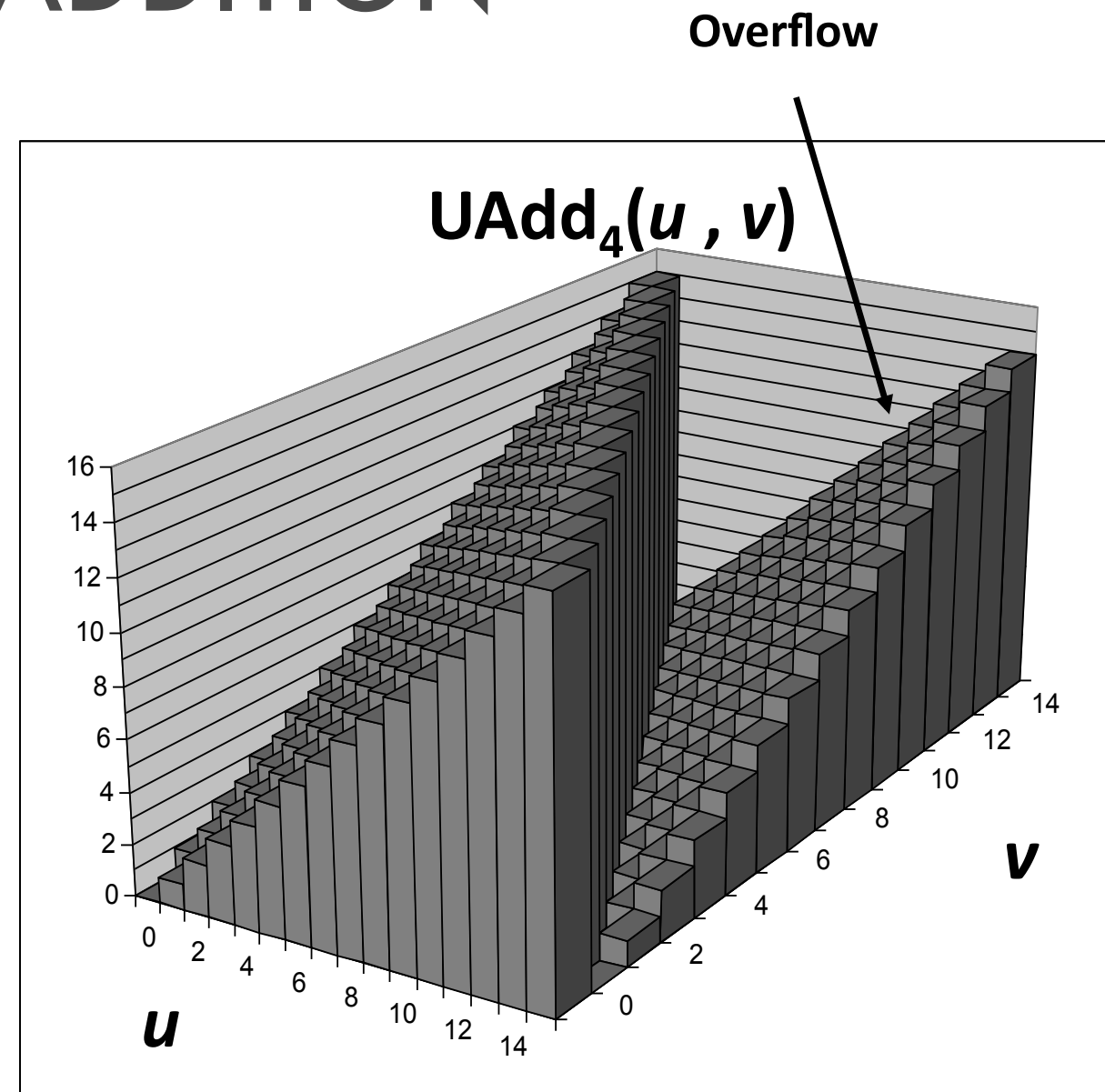
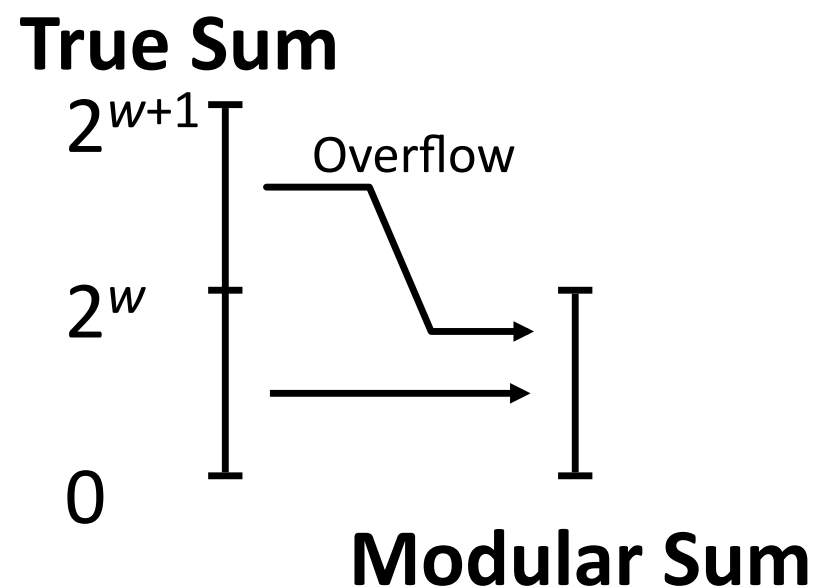
- Integer Addition
 - 4-bit integers u, v
 - Compute true sum $\text{Add}_4(u, v)$
 - Values increase linearly with u and v
 - Forms planar surface



INTEGER ARITHMETIC

VISUALIZING UNSIGNED ADDITION

- Wraps Around
 - If true sum $\geq 2^w$
 - At most once



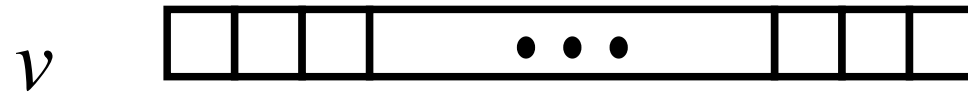
INTEGER ARITHMETIC

TWO'S COMPLEMENT ADDITION

Operands: w bits



+



True Sum: $w+1$ bits



Discard Carry: w bits

$\text{TAdd}_w(u, v)$



- TAdd and UAdd have Identical Bit-Level Behavior
 - Signed vs. unsigned addition in C:
 - `int s, t, u, v;`
 - `s = (int) ((unsigned) u + (unsigned) v);`
 - `t = u + v`
 - Will give `s == t`



INTEGER ARITHMETIC

TWO'S COMPLEMENT ADDITION

Operands: w bits



+

v



True Sum: $w+1$ bits

$u + v$



Discard Carry: w bits

$\text{TAdd}_w(u, v)$



- TAdd and UAdd have Identical Bit-Level Behavior
- Signed vs. unsigned addition in C:
 - `int s, t, u, v;`
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 - `t = u + v;`
- Will give `s == t`



INTEGER ARITHMETIC

TWO'S COMPLEMENT ADDITION

Operands: w bits

u 

+

v



True Sum: $w+1$ bits

$u + v$



Discard Carry: w bits

$\text{TAdd}(u, v)$



Why don't you try a few to convince yourself it's true?

- TAdd and UAdd have
- Signed vs. unsigned addition in C:
 - `int s, t, u, v;`
 - `s = (int) ((unsigned) u + (unsigned) v);`
 - `t = u + v;`
- Will give `s == t`

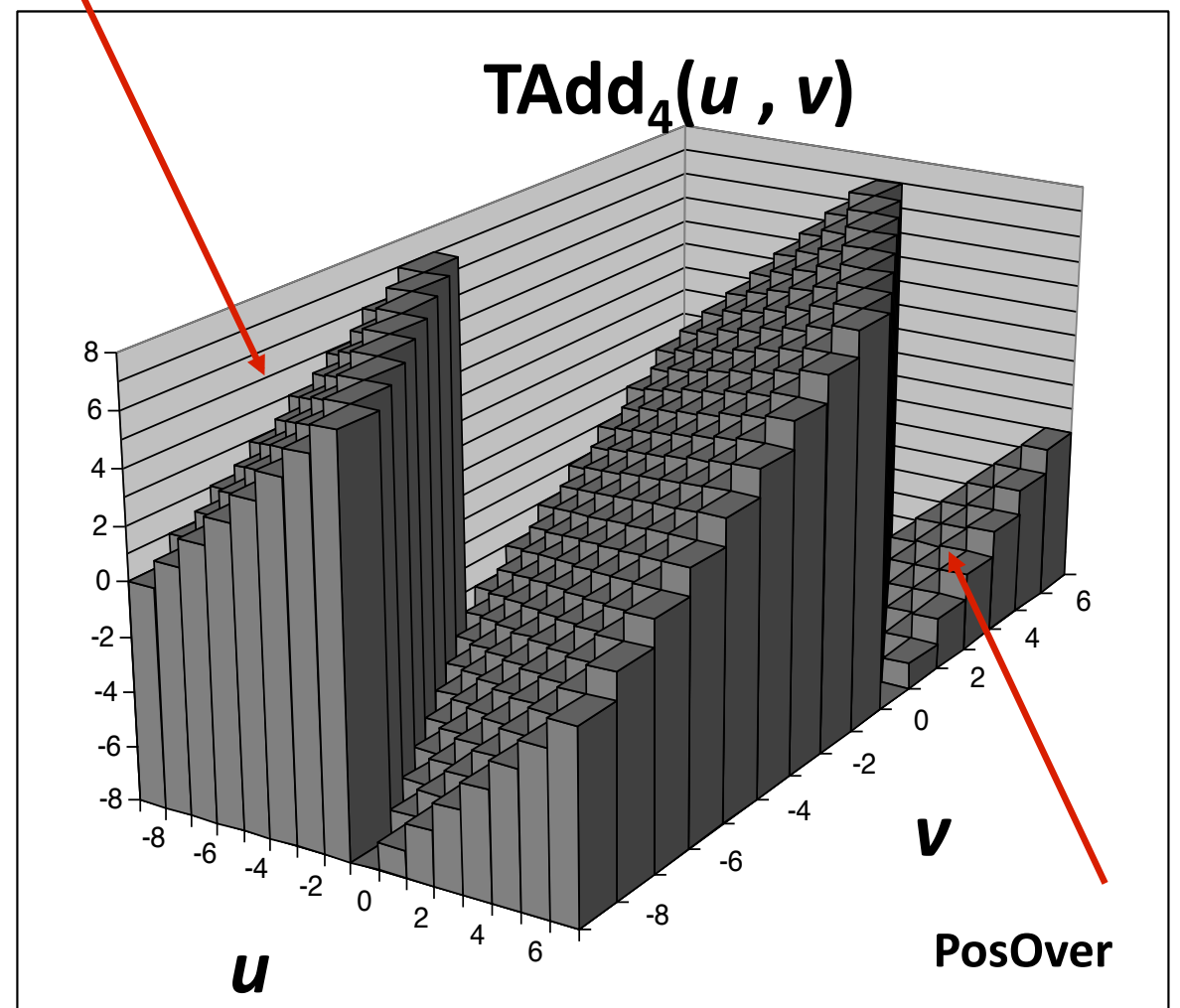


INTEGER ARITHMETIC

VISUALIZING 2'S COMPLEMENT ADDITION

- Values
 - 4-bit two's comp.
 - Range from -8 to +7
- Wraps Around
 - If $\text{sum} \geq 2^{w-1}$
 - Becomes negative
 - At most once
 - If $\text{sum} < -2^{w-1}$
 - Becomes positive
 - At most once

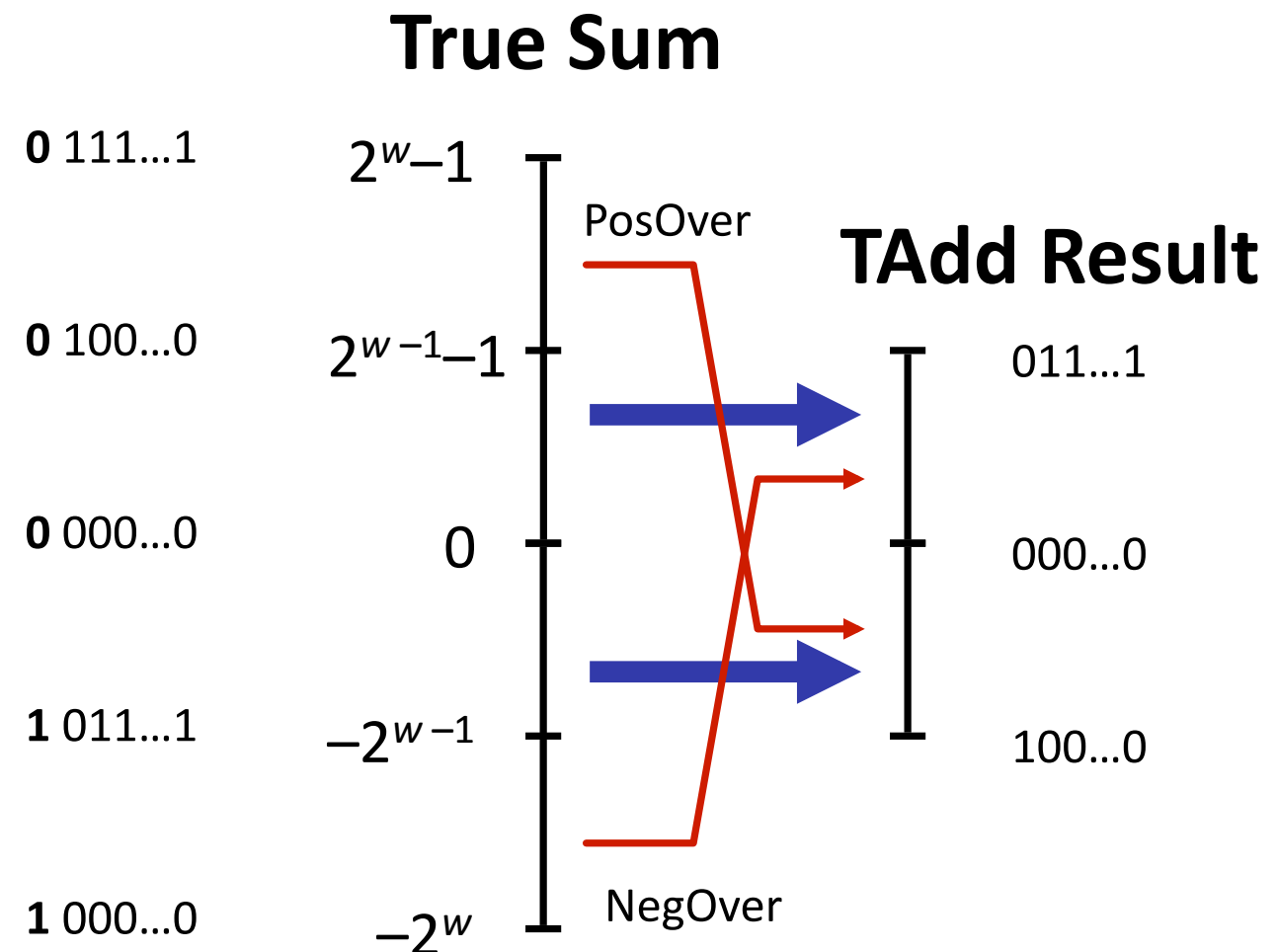
NegOver



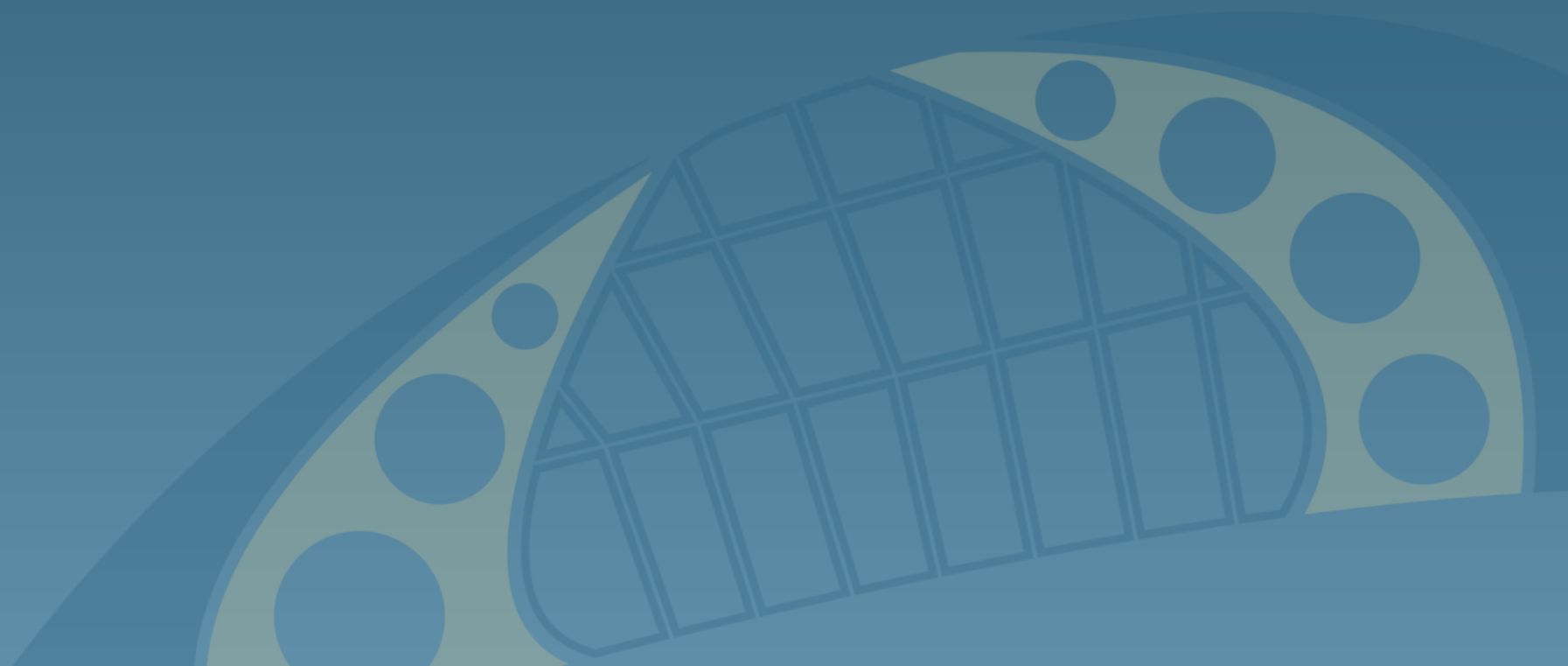
INTEGER ARITHMETIC

TADD OVERFLOW

- Functionality
 - True sum requires $w+1$ bits
 - Drop off MSB
 - Treat remaining bits as 2's comp. integer



INTEGER MULTIPLICATION



INTEGER ARITHMETIC

MULTIPLICATION

- Goal: Computing Product of w -bit numbers x, y
 - Either signed or unsigned
- But, exact results can be bigger than w bits
 - Unsigned: up to $2w$ bits
 - Result range: $0 \leq x * y \leq (2^w - 1) * (2^w - 1) = 2^{2w} - 2^{w+1} + 1$
- Two's complement min (negative): Up to $2w-1$ bits
 - Result range: $x * y \geq (-2^{w-1}) * (2^{w-1} - 1) = -2^{2w-2} + 2^{w-1}$
 - Two's complement max (positive): Up to $2w$ bits, but only for $(TMin_w)^2$
 - Result range: $x * y \leq (-2^{w-1}) * (-2^{w-1}) = 2^{2w-2}$



INTEGER ARITHMETIC

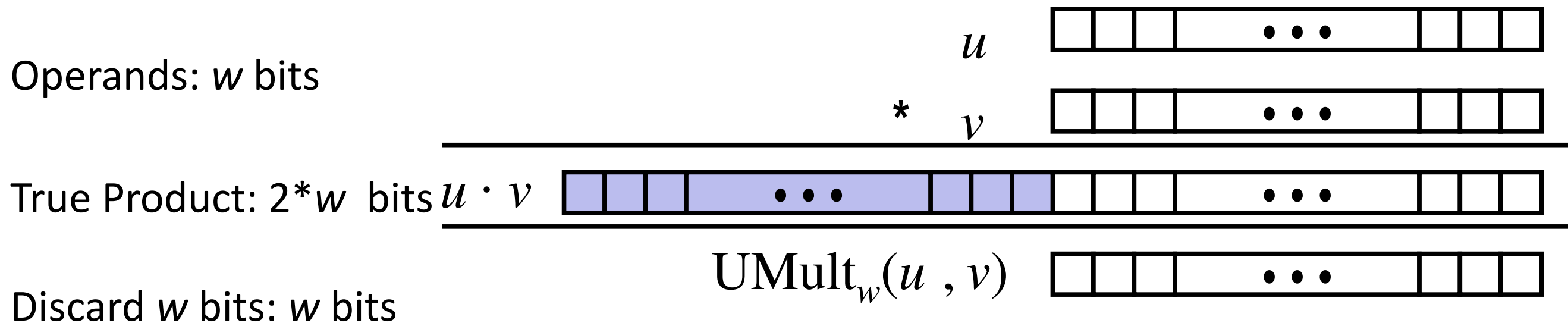
MULTIPLICATION

- So, maintaining exact results...
 - would need to keep expanding word size with each product computed
 - is done in software, if needed
 - e.g., by “arbitrary precision” arithmetic packages



INTEGER ARITHMETIC

UNSIGNED MULTIPLICATION IN C



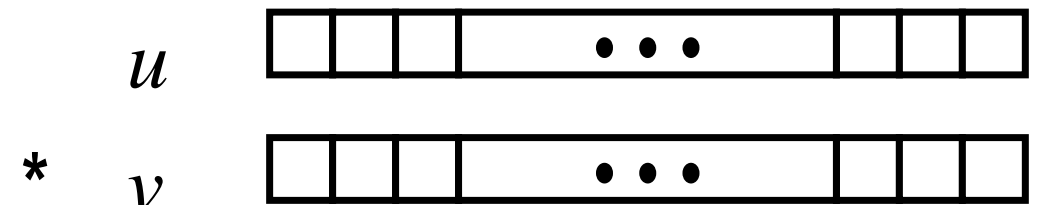
- Standard Multiplication Function
 - Ignores high order w bits
- Implements Modular Arithmetic
 - $\text{UMult}_w(u, v) = u \cdot v \bmod 2^w$



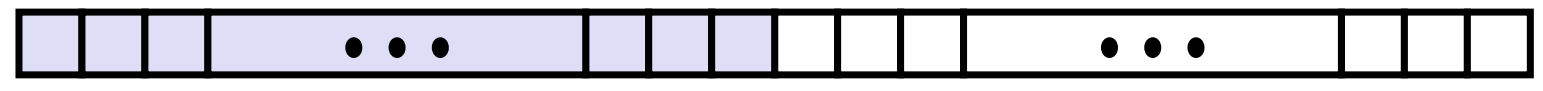
INTEGER ARITHMETIC

SIGNED MULTIPLICATION IN C

Operands: w bits



True Product: $2*w$ bits $u \cdot v$



Discard w bits: w bits

$\text{TMult}_w(u, v)$



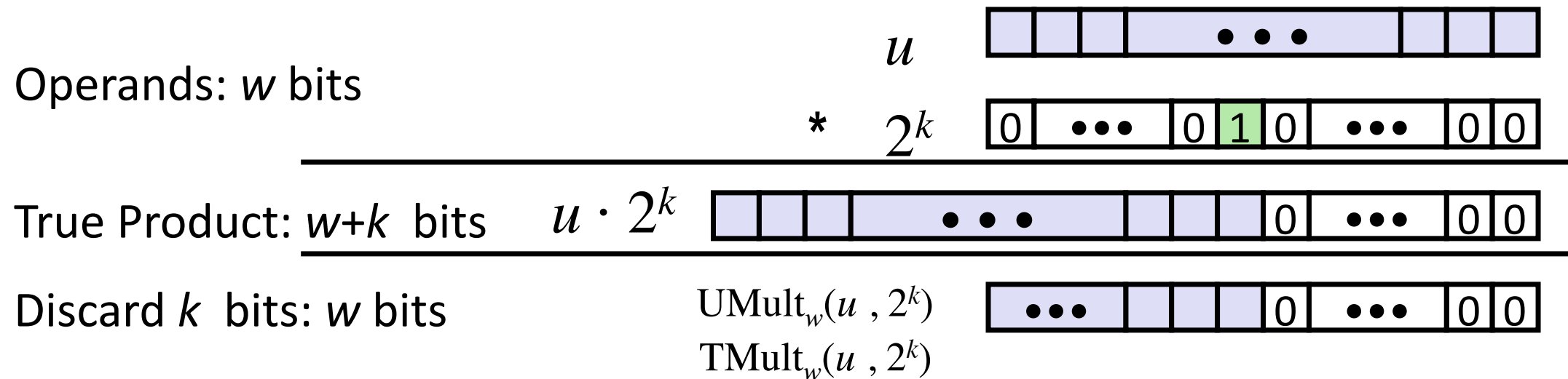
- Standard Multiplication Function
 - Ignores high order w bits
 - Some of which are different for signed vs. unsigned multiplication
 - Lower bits are the same



INTEGER ARITHMETIC

POWER-OF-2 MULTIPLY WITH SHIFT

- Operation
 - $u \ll k$ gives $u * 2^k$
 - Both signed and unsigned

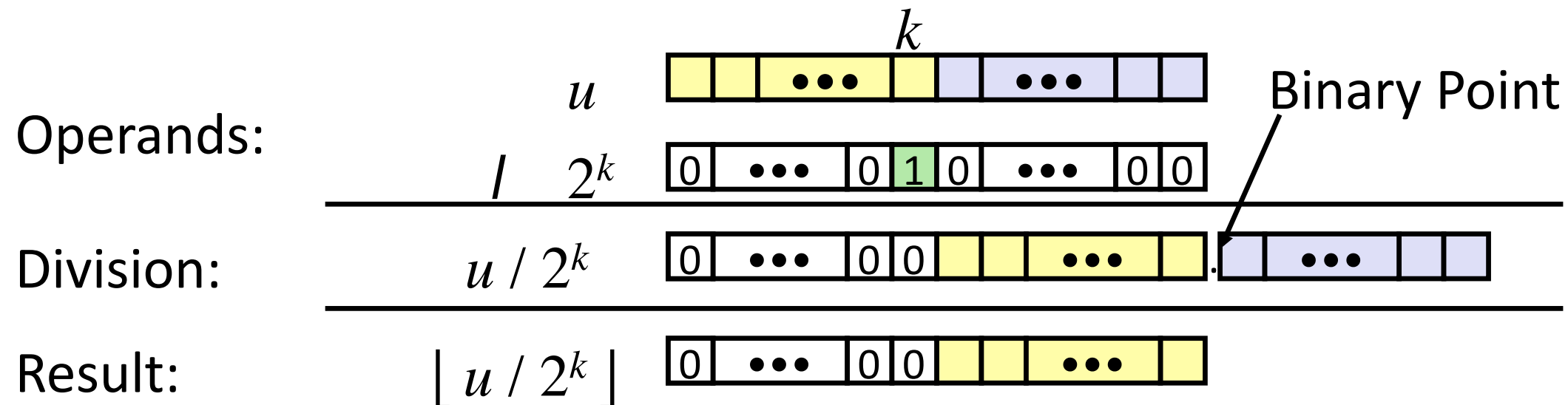


- Examples
 - $u \ll 3 == u * 8$
 - $(u \ll 5) - (u \ll 3) == u * 24$
- Most machines shift and add faster than multiply
 - Compiler generates this code automatically



INTEGER ARITHMETIC

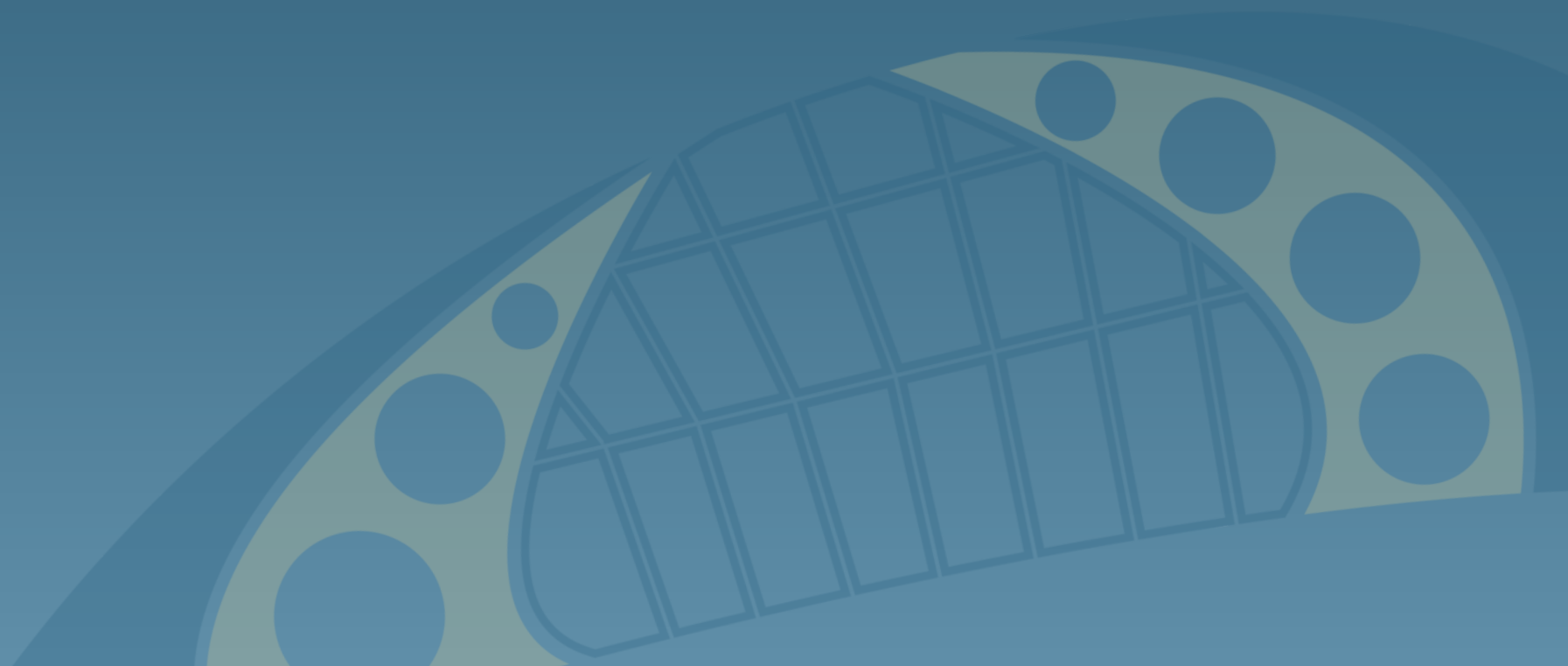
UNSIGNED POWER-OF-2 DIVIDE WITH SHIFT



	Division	Computed	Hex	Binary
x	15213	15213	3B 6D	00111011 01101101
x >> 1	7606.5	7606	1D B6	00011101 10110110
x >> 4	950.8125	950	03 B6	00000011 10110110
x >> 8	59.4257813	59	00 3B	00000000 00111011

- Quotient of Unsigned by Power of 2
 - $u \gg k$ gives $\text{floor}(u / 2^k)$
 - Uses logical shift

INTEGER ARITHMETIC SUMMARY



INTEGER ARITHMETIC

ARITHMETIC: BASIC RULES

- Addition:
 - Unsigned/signed: Normal addition followed by truncate, same operation on bit level
 - Unsigned: addition mod 2^w
 - Mathematical addition + possible subtraction of 2^w
 - Signed: modified addition mod 2^w (result in proper range)
 - Mathematical addition + possible addition or subtraction of 2^w



INTEGER ARITHMETIC

ARITHMETIC: BASIC RULES

- Multiplication:
 - Unsigned/signed: Normal multiplication followed by truncate, same operation on bit level
 - Unsigned: multiplication mod 2^w
 - Signed: modified multiplication mod 2^w (result in proper range)



INTEGER ARITHMETIC

WHY SHOULD I USE UNSIGNED?

- Don't use without understanding implications
 - Easy to make mistakes
 - `unsigned i;`
 - `for (i = cnt-2; i >= 0; i--)`
 - `a[i] += a[i+1];`
 - Can be very subtle
 - `#define DELTA sizeof(int)`
 - `int i;`
 - `for (i = CNT; i-DELTA >= 0; i-= DELTA)`
 - `. . .`



INTEGER ARITHMETIC

COUNTING DOWN WITH UNSIGNED

- Proper way to use unsigned as loop index
 - `unsigned i;`
 - `for (i = cnt-2; i < cnt; i--)`
 - `a[i] += a[i+1];`
- See Robert Seacord, *Secure Coding in C and C++*
 - C Standard guarantees that unsigned addition will behave like modular arithmetic
 - $0 - 1 \rightarrow \text{UMax}$



INTEGER ARITHMETIC

COUNTING DOWN WITH UNSIGNED

- Even better
 - `size_t i;`
 - `for (i = cnt-2; i < cnt; i--)`
 - `a[i] += a[i+1];`
 - Data type `size_t` defined as unsigned value with length = word size
 - Code will work even if `cnt = UMax`
 - What if `cnt` is signed and `< 0`?



INTEGER ARITHMETIC

WHY SHOULD I USE UNSIGNED? (CONT.)

- Do Use When Performing Modular Arithmetic
 - Multiprecision arithmetic
- Do Use When Using Bits to Represent Sets
 - Logical right shift, no sign extension



WORK IT OUT

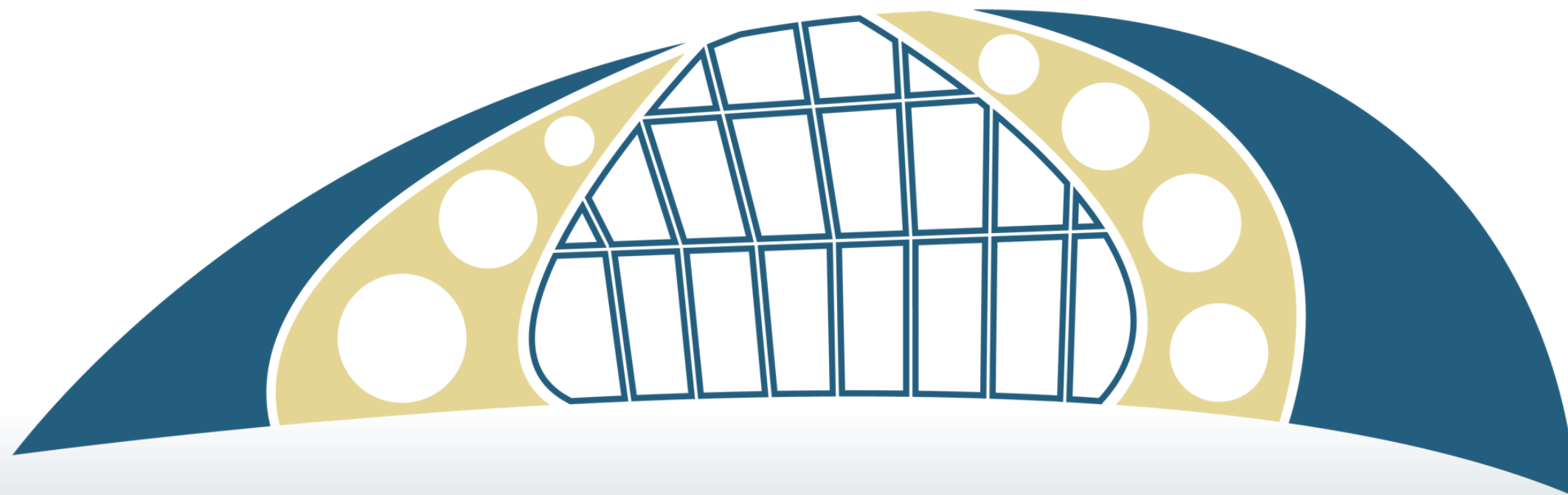
WHAT ARE M AND N?

```
#define M      /* Mystery number 1 */
#define N      /* Mystery number 2 */
int arith(int x, int y) {
    int result = 0;
    result = x*M + y/N; /* M and N are mystery numbers. */
    return result;
}
```

Original code

Optimized Code

```
/* Translation of assembly code for arith */
int optarith(int x, int y) {
    int t = x;
    x <<= 5;
    x -= t;
    if (y < 0) y += 7;
    y >>= 3; /* Arithmetic shift */
    return x+y;
}
```



WESTMONT **INSPIRED**
— COMPUTING LAB —